

Computer Networks CS3611

Transport Layer-Part 1

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The slides are adapted from those provided by Prof. Romit Roy Choudhury.

Chapter 3: Transport Layer

Our goals:

- understand principles
 behind transport layer
 services:
 - multiplexing/demultiplexing
 - reliable data transfer
 - flow control
 - congestion control

- learn about transport layer protocols in the Internet:
 - UDP: connectionless transport
 - TCP: connection-oriented transport
 - TCP congestion control

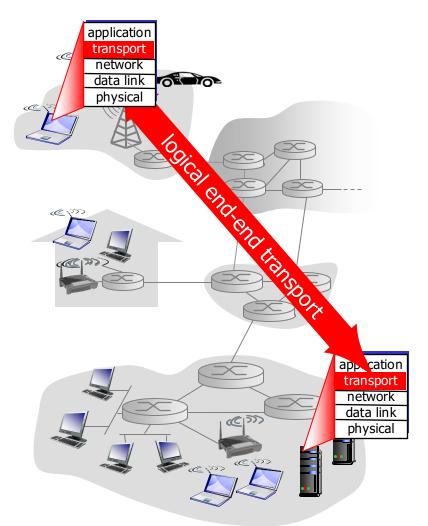
Chapter 3 outline

- 3.1 Transport-layer services
- 3.2 Multiplexing and demultiplexing
- 3.3 Connectionless transport: UDP
- 3.4 Principles of reliable data transfer

- 3.5 Connection-oriented transport: TCP
 - segment structure
 - reliable data transfer
 - flow control
 - connection management
- 3.6 Principles of congestion control
- **3.7 TCP congestion control**

Transport services and protocols

- provide *logical communication* between app processes running on different hosts
- □ transport protocols run in end systems
 - sender: breaks app messages into segments, passes to network layer
 - receiver: reassembles segments into messages, passes to app layer
- more than one transport protocol available to apps
 - Internet: TCP and UDP



Transport vs. network layer

- network layer: logical communication between hosts
- transport layer: logical communication between processes
 - relies on, enhances, network layer services

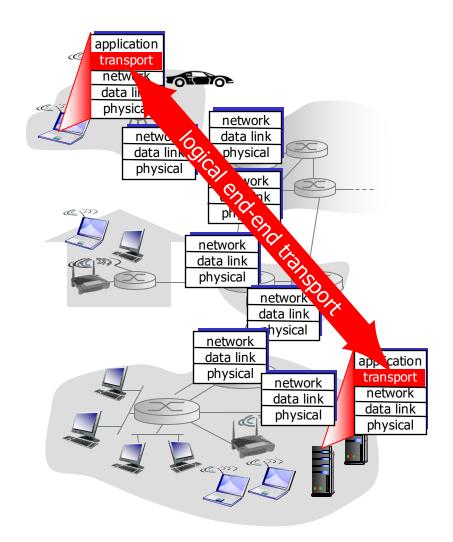
Household analogy:

12 kids sending letters to 12 kids

- \Box processes = kids
- app messages = letters in envelopes
- \Box hosts = houses
- □ transport protocol = Ann to Bill
- network-layer protocol = postal service

Internet transport-layer protocols

- reliable, in-order delivery (TCP)
 - congestion control
 - flow control
 - connection setup
- unreliable, unordered delivery: UDP
 - no-frills extension of "besteffort" IP
- services not available:
 - delay guarantees
 - bandwidth guarantees

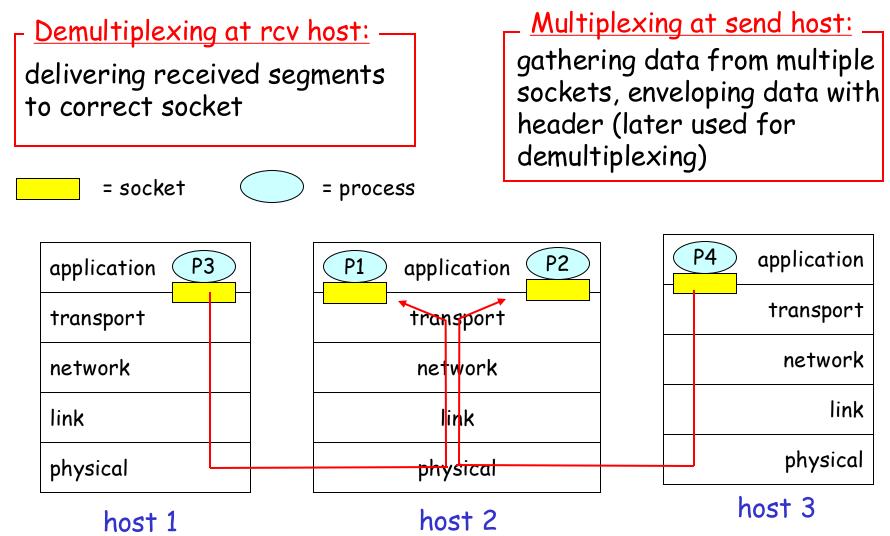


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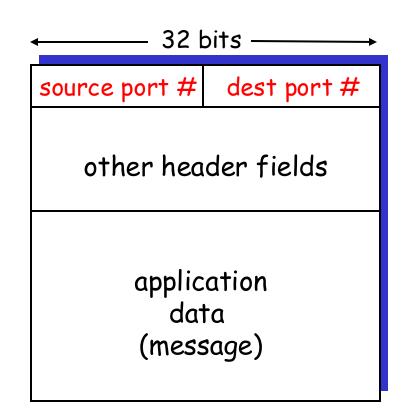
Multiplexing/demultiplexing



How demultiplexing works

host receives IP datagrams

- each datagram has source IP address, destination IP address
- each datagram carries 1 transportlayer segment
- each segment has source, destination port number
- host uses IP addresses & port numbers to direct segment to appropriate socket



TCP/UDP segment format

Connectionless demultiplexing

UDP socket identified by twotuple:

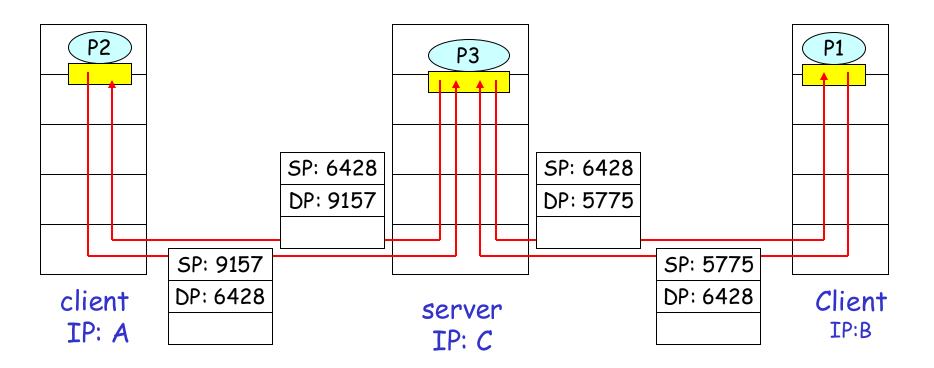
(dest IP address, dest port number)

When host receives UDP segment:

- checks destination port number in segment
- directs UDP segment to socket with that port number

 IP datagrams with different source IP addresses and/or source port numbers directed to same socket

Connectionless demux (cont)



SP provides "return address"

Connection-oriented demux

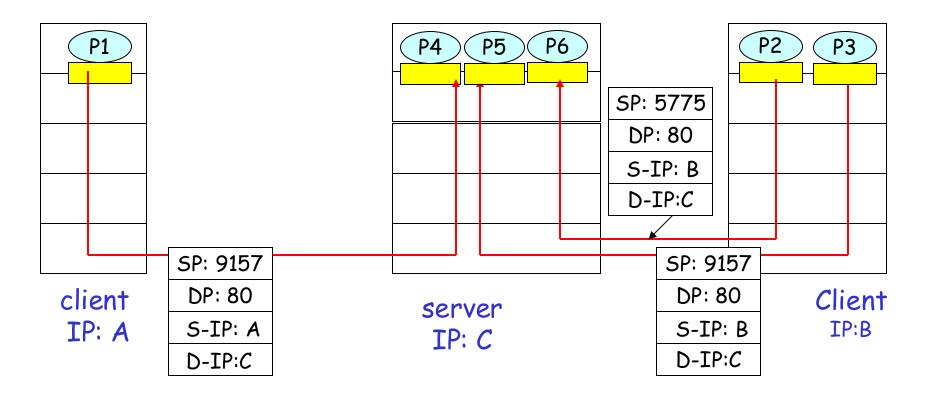
- TCP socket identified by 4tuple:
 - source IP address
 - source port number
 - dest IP address
 - dest port number
- recv host uses all four values to direct segment to appropriate socket

- Server host may support many simultaneous TCP sockets:
 - each socket identified by its own 4-tuple
- Web servers have different sockets for each connecting client
 - non-persistent HTTP will have different socket for each request

Connection-oriented demux (cont)

= socket

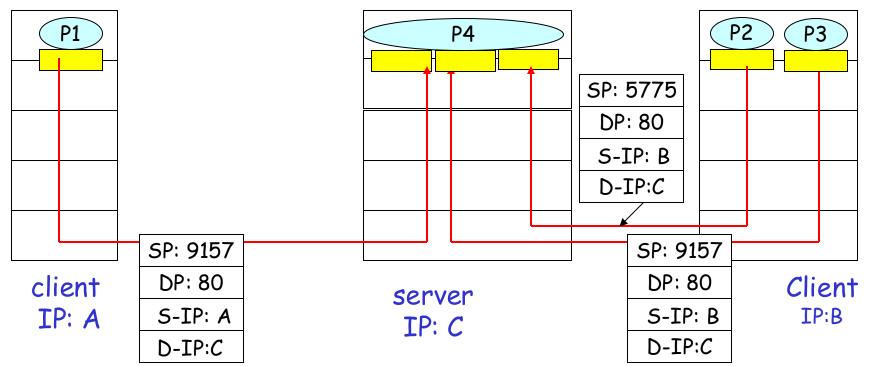
= process



<u>Connection-oriented demux:</u> <u>Threaded Web Server</u>

= socket





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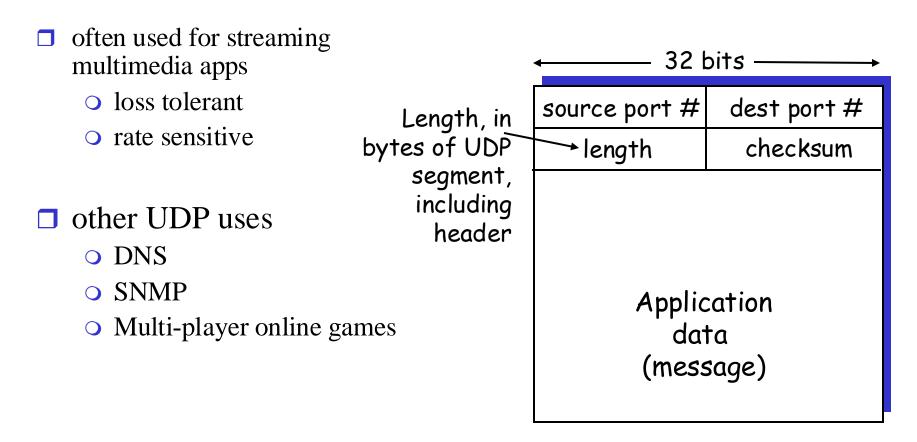
UDP: User Datagram Protocol [RFC 768]

- "no frills," "bare bones" Internet transport protocol
- "best effort" service, UDP segments may be:
 - lost
 - delivered out of order to app
- connectionless:
 - no handshaking between UDP sender, receiver
 - each UDP segment handled independently of others

Why is there a UDP?

- no connection establishment (which can add delay)
- simple: no connection state at sender, receiver
- small segment header
- no congestion control: UDP can blast away as fast as desired

UDP: more



UDP segment format

UDP checksum

Goal: detect "errors" (e.g., flipped bits) in transmitted segment

Sender:

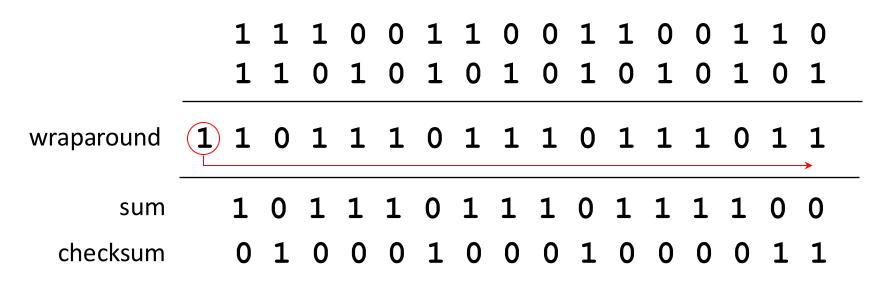
- treat segment contents as sequence of 16-bit integers
- checksum: addition (1's complement sum) of segment contents
- sender puts checksum value into UDP checksum field

Receiver:

- compute checksum of received segment
- check if computed checksum equals checksum field value:
 - NO error detected
 - YES no error detected. *But maybe errors nonetheless?* More later

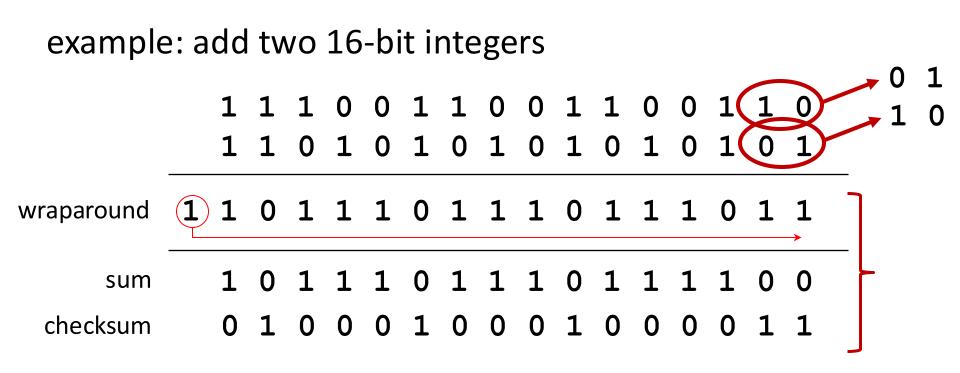
UDP checksum

example: add two 16-bit integers



Note: when adding numbers, a carryout from the most significant bit needs to be added to the result

UDP checksum: weak protection



Even though numbers have changed (bit flips), *no* change in checksum!